Firming Renewable Power with Demand Response: An End-to-end Aggregator Business Model

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Abstract

Environmental concerns have spurred greater reliance on variable renewable energy resources (VERs) in electric generation. Under current incentive schemes, the uncertainty and intermittency of these resources impose costs on the grid, which are typically socialized across the whole system, rather than born by their creators. We consider an institutional framework in which VERs face market imbalance prices, giving them an incentive to produce higher-value energy subject to less adverse uncertainty. In this setting, we consider an “aggregator” that owns the production rights to a VER’s output, and also signs contracts with a population of demand response (DR) participants for the right to curtail them in real time, according to a contractually specified probability distribution. The aggregator bids a day ahead offer into the wholesale market, and is able to offset imbalances between the cleared day-ahead bid and the realized VER production by curtailing DR participants’ consumption according to the signed contracts. We consider the optimization of the aggregator’s end-to-end problem: designing the menu of DR service contracts using contract theory, bidding into the wholesale market, and dispatching DR consistently with the contractual agreements. We do this in a setting in which wholesale market prices, VER output, and participant demand are all stochastic, and possibly correlated.

Keywords: Electricity Markets, Demand Response, Aggregator, Business Model, Renewables Integration, Market Design, Screening Mechanisms

JEL Classification: D11, D45, D47, Q41, Q42
1 Introduction

1.1 Background and Motivation

Environmental concerns regarding global warming and the adverse health effects of emissions produced by fossil fuel generation have led to a greater reliance on renewable sources of generation, such as solar and wind, which are inherently variable and uncertain. This trend is accompanied by increased proliferation of distributed resources, storage, and smart grid technologies for metering and control, which facilitate demand response and greater observability of the grid. As a result, the electric power industry faces new challenges in planning and operation of the power system that require new institutional and regulatory frameworks, along with appropriate market mechanisms to achieve productive and allocative efficiencies.

While the conventional approach to mitigating adverse uncertainty and variability on the supply and demand sides has been increased reliance on reserves and flexible generation units, this approach is expensive, and will undermine the economic and environmental goals of renewables integration. Mobilizing demand side flexibility enabled by smart metering and other smart grid technologies to mitigate the uncertainty and variability of renewable resources is a sustainable solution for addressing the operational challenges posed by massive integration of renewables.

Alternative approaches to integrate renewable resources into the power grid and facilitate demand response have been proposed and experimented with by policy makers around the world, and have been the subject of numerous academic studies in the economics and power system literature. From an economics perspective, the gold standard approach to achieving production and allocative efficiency is a centralized market where all renewable resources and conventional resources are pooled together with demand side resources, responding to real time marginal prices set through a market clearing mechanism. However, while such an approach may serve as a useful benchmark, it is impractical, as it would require the system operator to collect information and co-optimize the dispatch of a vast number of resources including conventional generation, renewables and participating demand side resources (PDR). The computational and institutional barriers to such a centralized approach calls for more pragmatic second-best alternatives with more manageable scope. Recent regulatory initiatives such as “Reforming the Energy Vision” (REV) initiated by the New York Public Service Commission (PCS) promote a more decentralized approach as a way to facilitate the integration of decentralized renewable resources and demand response (MDPT Working Group, 2015). Likewise, the concept of aggregators that can pool demand side resources and act as intermediaries, offering load reduction into the wholesale market, has been popularized by the emergence of commercial entities such as EnerNOC. The scope of such aggregation can be expanded to include behind-the-meter resources and distributed renewable resources.

In this paper we propose and analyze an aggregator business model that assembles a portfolio of variable energy resources (VER) such as wind, and of flexible demand response (DR), with the purpose of producing a firm and controllable bundled energy resource that can be offered into
the ISO wholesale day ahead market. We presume that the aggregator is in a position to acquire detailed information and enter into contractual arrangements that would enable it to mobilize the DR flexibility so as to offset the VER uncertainty and variability. Such a bundled resource will relieve the ISO from having to procure additional reserves or other ancillary service products for the purpose of mitigating renewables intermittency.

Our premise in this paper is that future regulatory reforms will provide incentives to VER to firm up their output and induce loads to surrender their flexibility. On the VER side, such incentives will be enabled when subsidies to renewables such as feed in tariffs will be replaced by nondiscriminatory market mechanisms. Under such a mechanism, uncertain resources would bear the cost they impute on the system, whereas flexible resources are rewarded for the flexibility. Furthermore, VER will have to schedule their forecasted production and be subject to deviation settlements in the real time market like other resources, whereas firmed up VER will be eligible for capacity payments through resource adequacy mechanisms. On the demand side, ex ante contractual agreements with an aggregator that compensate the customer for forgone consumption and “information rents” should provide incentives for load to reveal and trade their flexibility.

The two principal forms of demand response are direct load control, wherein the aggregator physically constrains participants’ consumption during scarcity events, and price-based control, wherein the participants face real-time prices that reflect current system conditions.\(^1\) Direct load control has been studied in theory (Chao, 1983) and implemented in practice, particularly in contexts such as air conditioner cycling (RLW Analytics, 2007). It has the advantage in terms of system reliability, because the response is more predictable; as well as with respect to billing simplicity and predictability, because the customer does not face state-dependent prices. On the other hand, price-based control provides customers with more flexibility. According to standard microeconomic models, the most economically efficient form of control is real-time pricing, because it ensures that customers consume exactly when their marginal benefit is greater than the instantaneous marginal cost of power production (Borenstein, 2005; Caramanis et al., 1983; Holland and Mansur, 2006).\(^2\) If the consumer’s demand curve for power were constant over time, then a direct load control contract linked to spot prices would result in the same consumption decisions as real time pricing (Chao and Wilson, 1987).

Restructured electricity markets are premised on treating electricity at the wholesale level as a homogeneous commodity that is produced and traded based on fluctuating price signals. We argue, however, that at the retail level electricity can be offered as a quality differentiated service with predetermined prices and uncertain availability (quantity control). Such uncertainty is realized through direct load control, or customer response to a load control signal subject to a noncompliance penalty. The above perspective, which has been articulated by Oren (2013), is the underlying

\(^1\)Caramanis, Schweppe and Tabors categorized direct load control as “price / quantity transaction,” and price-based control as “price only transactions” (Caramanis et al., 1983).

\(^2\)The standard analysis ignores intertemporal interactions; but see, for example, Tsitsiklis and Xu (2015) for an extension to pricing for contribution to system-wide ramping cost.
paradigm explored in this paper and we will not attempt to contrast it
with a real time pricing approach, which as we concede, represents the
“economic gold standard.” Specifically, we consider a profit-maximizing
aggregator contracting ex ante with DR participants for the right to send
curtailment signal with a specified probability (or, more generally, in
specified states of the world, as reflected by a publicly observable index).
The curtailment signal effectively raises the participant’s price for calling
energy from a particular capacity increment from its original retail rate
$R$, to an exogenously determined “penalty price,” $H > R$. That is, the
capacity increment is an option, and curtailment raises the strike price.
The case where $H = \infty$ can be interpreted as direct load control. This
generalizes plans like PG&E’s SmartRate plan, which raises the customer’s
tariff for 15 days a year or less (PG&E, 2015). In our generalization,
different slices of the household’s consumption capacity have different
probabilities of facing curtailment/penalty rate signal. Combined with a
model of stochastic valuations for service, this approach models two kinds
of imperfect or fractional DR yield: DR that fails to materialize because the
customer would not have consumed in the first place (the ex post valuation
of consumption is less than $R$), and DR that fails to materialize because
the customer’s ex post valuation is higher than the penalty price, $H$. In
either case we assume that the valuations are constant throughout the time
interval, and each valuation is for the energy from an infinitesimal capacity
slice, so we do not consider the possibility of partial exercise of a capacity
increment within a period. However, by “stacking” these increments, the
model generalizes to horizontal load slices that can be fractionally utilized,
at a constant level during the period. Less-than-infinite penalties may be
a happy medium between the intrusiveness of a hard constraint, and the
complexity of a real-time price.

Our proposed business model is based on a “fuse-control paradigm”
(Margellos and Oren, 2015) when the aggregator manages the service
quality for the aggregate consumption by imposing a capacity constraint, or
by signaling a capacity threshold above which the penalty will be imposed,
and leaves the decision of allocating the available power to devices behind
the meter to the household. This is a less intrusive alternative to direct
curtailment of individual devices, such as air-conditioner cycling programs,
for instance. Furthermore, delegating the behind-the-meter allocation
allows the customer to reflect intertemporal variations in preferences
for different electricity uses, and capture the effect of behind-the-meter
variable resources such as solar panels, local storage devices, and deferrable
energy uses such as electric vehicle charging, HVAC etc. In our model the
aggregator is assumed to submit price-contingent hourly offers into the ISO
day ahead market and dispatch curtailment signals to its contracted load
based on the awarded quantities in the day ahead market, the realized
renewable output, and the deviation settlement prices. Our end-to-end
approach seeks to co-optimize the contract design on the demand side
with the aggregator’s bidding strategy in the ISO day ahead market and
the DR deployment strategy.
1.2 Related work

Motivated by the same concern about the subsidization of VERs’ contribution to reserve costs, Bitar et al. (2011, 2012) consider several stylized market models for renewable power. They use a newsvendor-type model to quantify the effect of imbalance charges on the offer behavior and profit of a renewable producer, and to quantify, for example, the value of forecast improvement in this policy environment. Their second model is a market for reliability-differentiated power, originally studied by Tan and Varaiya (Tan, 1990; Tan and Varaiya, 1993). In this model, the producer owns a stochastic power resource, and sells its entire production in advance without using reserves, by offering contracts with imperfect service reliability. Our model can be seen as a synthesis and generalization of the two models above. Tan (1990) has established, in a general equilibrium analysis, the existence of a welfare-maximizing price-taking equilibrium, and certain structural characteristics thereof, in his market model with finitely many customer types. In contrast, our setting is monopoly / monopoly profit maximization with a continuum of consumer types, and each unit of demand has imperfect knowledge of its valuation at the time of contracting.

We cast the problem of designing an optimal menu of variable-reliability demand response contracts as a variation on the classic monopoly screening problem from contract theory. Our approach to embedding this screening problem in a wholesale electricity market follows the literature on priority service, particularly Chao and Wilson (1987) and Chao (2012). However, that literature has focused on perfect competition or regulated social welfare maximization, and abstracts away from the scheduling and recourse decisions of individual producers. Because we are interested in new business models that manage imbalance, we update the priority service approach in a profit maximization setting, where imbalance cost is reflected by imbalance prices.

Another point of contrast with Chao (2012) is that our stochastic demand model disaggregates the aggregate demand curve along the quantity axis, and then adds post-contracting noise to valuations, in a manner similar to Courty and Li’s (Courty and Li, 2000) sequential screening model. However, in contrast to most screening environments, including that of Courty and Li (2000), our producer’s contracting problem is embedded in a newsvendor-like problem, with asymmetric linear prices for positive and negative imbalance. As a consequence, the aggregator’s benefit is not linear (i.e. is not an expectation) over a type distribution. The aggregator-cum-producer co-optimizes its demand response menu with a day ahead offer quantity, with the demand response providing recourse in case of real-time imbalance.

Recently, Crampes and Léautier (2015) have used contract theory to study the welfare effects of allowing demand response participation in adjustment markets, when DR participants have private information about their utility from consumption. In their setting, vertically integrated pro-

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3But see Wilson (1993), who treats profit maximization but in a slightly different setting from ours.
producers contract as retailers with consumers in the first stage, so that the consumers “buy the baseline” via monopoly contracting, and producers incur the obligation to produce the contracted amount. Then, in the second stage, all producers experience an identical supply shock (capacity failure), and both producers and consumers can participate in a competitive adjustment market. They employ a stylized, two-type model to show that there exist cases in which allowing consumers to participate in an adjustment market reduces social welfare, by creating sufficiently large distortions in the retail contract market. Their model is suited to showing the existence of certain distortions that regulators should be wary of. In contrast, we model the market demand curve as comprising a continuum of types. This provides a more detailed, less stylized account of how an aggregator should optimize a production offer and DR dispatch policy with knowledge of market statistics, renewable output, demand conditions, etc.

1.3 Summary of Results

We analyze the profit-maximizing policy of an aggregator, which owns a renewable resource and contractually purchases DR as a monopsonist, in order to sell a bundled product into a competitive wholesale market with linear imbalance prices. In our setting, the aggregator purchases less than the socially efficient level of demand response. This follows from two assumptions: (1) the aggregator is a monopsonist; and (2) DR participants already own unlimited options to consume. We discuss remedies for this, from a welfare maximization perspective, in the final discussion.

1.4 Introduction to the model

We consider the profit maximization problem of an aggregator. This aggregator has two sources from which it produces energy: a VER (“wind”) with known probability distribution over production quantities, and a population of DR participants, with whom it signs contracts ex ante (say, at the beginning of the season) giving the aggregator the right to curtail them with specified probabilities. The market system operator treats reductions in participants’ consumption, induced by curtailment, as the aggregator’s production. The DR participants have private information regarding their valuation for service, and the aggregator acts as a monopsonist purchaser of rights to curtail increments of their capacity with specified probabilities. We analyze the aggregator’s problem as a “screening problem” (Börgers, 2010) in which the aggregator’s benefit function reflects its participation in the wholesale electricity market, as we describe presently.

The aggregator bundles the VER and DR production for sale into a wholesale electricity market. It chooses an energy offer policy \( q \) into the day-ahead (DA) market, contingent on DA information. If the DA offer is made contingent only on the price \( p \), then this price-contingent offer policy can be interpreted as a supply offer curve. Day-ahead, the aggregator receives revenue \( p q \). In the real time dispatch (RT) stage, it learns the

\[ \text{However, Crampes and Léautier argue that the qualitative insights carry over to imperfectly competitive settings.} \]
wind outcome $s$, the prices $a$ and $b$ for positive and negative deviations respectively from the DA commitment quantity $q$, and chooses a set of DR participants to curtail. Ex post, this results in a net (“aggregated” or “bundled”) production quantity, $s + DR$. The aggregator then pays $b(q-s-DR)^+ - a(q-s-DR)^-$, $a < p < b$, to settle the difference between the ex post production and the DA commitment. The joint probability distribution over all information is known in advance. One might say that from the ex ante perspective, the aggregator’s problem is a probability distribution over newsvendor problems, and in each newsvendor problem, after the initial quantity choice, the aggregator can take recourse actions in response to observed “demand” (here we mean negative wind), by dispatching DR. The DR cost is nonlinear, determined by the economics of the screening problem. In general the probability distribution over DR actions may be constrained to conform to contracts negotiated ex ante, but this advance commitment has no economic effect: assuming the aggregator faces no statistical or computational limitations, DR can be treated purely as a form of real-time recourse (see Section 2.4.1), whose optimal distribution it can foresee and therefore commit to from the ex ante stage.

The population of demand response participants is modeled as a continuum of “increments” of capacity—i.e., potential consumption. Ignoring stochasticity of valuations, each increment is a differential, $dx$, on the quantity axis of the population aggregate demand curve. Each infinitesimal increment has private information, indexed by its type $\tau \in [\tau, \overline{\tau}] \subset \mathbb{R}$, parameterizing the distribution over its “ex post valuation” $\theta$ for power at the time of consumption. The types of increments are distributed according to a measure with associated distribution function $G$ and density $g$, with convex compact support $[\tau, \overline{\tau}] \subset \mathbb{R}$. The measure of a set of increments under $G$ represents the total potential consumption capacity of that set, in MW.

Before laying out the microeconomic model of how DR is produced and how much it costs, we can informally write the aggregator’s problem, from the ex ante perspective, as:

$$\max_{q, DR, T} J_{EA}(q, DR, T) = \max_{q, DR, T} E \left[ p \ q + a (DR + s - q)^+ - b (q - DR - s)^- \right] - T.$$  

(1)

Here $q$, $DR$, and $T$ are policy variables. The $DR$ dispatch is determined in real time, although in accordance with a policy determined ex ante, and the corresponding payment $T$ is made ex ante (although it could equivalently be made ex post, as discussed below). The exogenous random variables

\[\text{5Our model implies that the aggregator offers a menu of quality-differentiated service options (Mussa and Rosen, 1978) to each increment, independently of contracting with the rest. A household chooses a menu item for each of its increments, and its total utility and transfer are Lebesgue integrals over the corresponding utility and transfers for each increment. This treatment rules out quantity-based nonlinear price discrimination.} \]

\[\text{6The EA subscript indicates that this objective is an expectation from the ex ante perspective.} \]
are:

\[ p : \text{day ahead ("DA") price} \in [p_L, p_U] \]
\[ a : \text{overproduction payment rate} \in [a_L, a_U], \text{realized in RT} \]
\[ b : \text{shortfall penalty rate} \in [b_L, b_U], \text{realized in RT} \]
\[ s : \text{VER ("wind") realization}, \in [s_L, s_U], \text{realized in real time (RT)}, \]
\[ \{\theta^s_\tau : \tau \in [\tau_L, \tau_U]\} : \text{DR participants' valuations, realized ex post.} \]

The last item is a continuum of random variables: a process, although indexed by the type of the DR participant, rather than by time. It does not show up in the informal objective, but we will explain how it affects the DR quantity \( DR \), and payment \( T \).

1.5 Information and Decision Structure

In our most general analysis, random variables are realized, and decisions taken, in four temporal stages. As mentioned above, the contracting decision that may be made at the ex ante stage can be treated as if it is postponed until the RT stage. At each subsequent stage, information from the previous stage is retained, new payoff-relevant random variables are realized, and forecasts of the random variables in future stages may be updated. We denote the tuple of random variables at each time-stage, drawn from the set of possible events, as \( \omega \in \Omega \), with a subscript denoting that time-stage, and we denote the non-payoff-relevant component as \( \xi \) with the same subscript.

0. Ex ante (EA) stage. The aggregator learns all probability distributions. The aggregator offers the same menu of contracts to each member of the population of DR participants, and the DR participants select their preferred plans. The aggregator makes the aggregate payment \( T \) from Equation (1). \(^9\)

1. Day-ahead (DA) stage. The aggregator learns the DA price \( p : \omega_{DA} = (p, \xi_{DA}) \in \Omega_{DA} \), and chooses its offer quantity, \( q(\omega_{DA}) \). The function

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We use the same symbols for random variables and their realized values.

\(^8\)We assume that this process is jointly measurable with respect to the product measure induced by \( g \) and the probability measure over \( \theta \). Therefore the process admits an essentially unique decomposition into an idiosyncratic component and an aggregate component (Al-Najjar, 1995). We are not particularly interested in technical issues regarding measurability. Instead we simply posit, as suggested in Judd (1985), that the idiosyncratic noise obeys an exact strong law of large numbers. In a similar spirit, we make whatever assumptions necessary to license the application of Fubini’s theorem can be applied to exchange the order of integration with respect to \( g \) and expectation over the process \( \theta \), which should not be demanding, since each valuation takes a value in a compact interval \([\theta_L, \theta_U] \subset \mathbb{R}\). Further, the increments’ decision that is contingent on this information process involves no strategic interaction, so most of the potential technical complications do not arise.

\(^9\)This is a slight simplification: the collection of any penalties from increments, for violating curtailment signals, is also netted out from \( T \) ex post.
\( \omega_{\text{DA}} \rightarrow q(\omega_{\text{DA}}) \) can be interpreted as a supply function offered in the ISO DA market, if it is \( \mu \)-measurable.\(^{10}\)

2. Real time dispatch (RT) stage. The aggregator learns the imbalance prices \((a, b)\) and wind outcome \(s\): \( \omega_{\text{RT}} = \omega_{\text{DA}} \times (a, b, s, \xi_{\text{RT}}) \in \Omega_{\text{RT}}, \) and chooses the set of DR increments to send curtailment signals to: \( \{ \tau : k(\tau, \omega_{\text{RT}}) = 1 \} \).\(^{11}\) A general curtailment function is denoted as \( k : [\tau, \overline{\tau}] \times \Omega_{\text{RT}} \rightarrow [0, 1] \), where the value \( k(\tau, \omega_{\text{RT}}) \) is the ex ante probability that type \( \tau \) is curtailed in RT event \( \omega_{\text{RT}} \). In Assumption 2 below, we restrict attention to curtailment rules of the form \( k(\tau, \omega) = \mathbb{1}_{\{ \tau \leq \hat{\tau}(\omega_{\text{RT}}) \}} \).

3. Ex post (EP) stage. The participants’ valuations are realized: \( \omega_{\text{EP}} = \omega_{\text{RT}} \times \{ \theta_{\tau} : \tau \in [0, N] \} \); this determines the realized quantity of demand response. We denote the latter random variable (or its realization in event \( \omega_{\text{EP}} \)) as \( \text{DR}(\omega_{\text{EP}}, \hat{\tau}(\omega_{\text{RT}})) \).

The aggregator’s primary policy variables are thus

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q : \omega_{\text{DA}} \rightarrow q(\omega_{\text{DA}}) \in \mathbb{R} \\
k(\cdot, \cdot) : (\tau, \omega_{\text{RT}}) \rightarrow k(\tau, \omega_{\text{RT}}) \in [0, 1].
\]

We alternate between \( \hat{\tau} \) and \( k \) notation for the curtailment policy as convenient. The payment \( T \) is a decision, but the screening analysis lets us express the optimal \( T \) given a curtailment policy \( k \) as a functional of that policy.

This is a rather general description, which is suitable to our analysis of the DR contract design component of the aggregator’s problem. However, we only solve the aggregator’s whole problem (the “end-to-end problem”), which embeds the DR contracting into a wholesale market, in special cases. In these special cases, some of the information is realized at earlier stages than in the general case, or is never stochastic; that is, certain advance forecasts are assumed to be perfect.

We occasionally omit the DA, RT and EP subscripts on random outcomes when the referent is clear from the context.

1.6 Organization of remaining sections

The remainder of the paper proceeds as follows: In Section 2, we characterize the class of merit order curtailment policies and their corresponding contracts in our setting. This determines the cost of implementing a curtailment policy. In Section 3, we analyze the aggregator’s end-to-end problem, which embeds DR contracting and dispatch into a newsvendor-style wholesale market. In Section 3.1, we present the general model of the aggregator’s benefits from demand response. In Sections 3.2 and 3.3, we consider two special cases of the end-to-end problem that we can solve to successive degrees of explicitness. These give us insight into the structure of the aggregator’s end-to-end problem. Finally, we conclude and outline extensions and future directions.

\(^{10}\)In our example in Section 3.3, the DA forecasts of imbalance prices are assumed to be perfect, so they are effectively revealed in this stage as well.

\(^{11}\)The product notation “\( \times \)” denotes concatenation of ordered tuples.
In this section, we analyze the DR contracting process. In Section 2.1, we lay out the key economic assumptions that make our problem tractable. In Section 2.2, we introduce direct mechanisms and specify our capacity increments' utility function in a direct mechanism. In Section 2.3 we invoke the Revelation Principle, the foundation of mechanism design. Section 2.4.1, we note that contracting for merit order curtailment policies can also be implemented “in real-time,” at least if we treat the economic model literally, which gives us economic insight into the set of implementable contracts and policies.

2.1 A one-dimensional type space

The key feature of our analysis of the DR contractual screening problem is that we make two assumptions that are jointly sufficient to render the increments’ type space “one-dimensional.”

The first assumption is that conditional on any real-time outcome $\omega_{RT}$, “higher types” have a higher distribution over ex post valuations, in the sense of first-order stochastic dominance (FOSD). At the ex post stage, each capacity increment will consume if its realized valuation for consumption is sufficiently high. Since there is a continuum of increments, each one is infinitesimal. Therefore we can assume that no two increments have the same type, and that a distinct ex post valuation random variable is associated with each type.

**Definition 2.1 (Ex post valuation).** The “ex post valuation” is the dollar value that an increment of type $\tau$ derives from consumption in ex post event $\omega_{EP}$. It is a random variable, with value

$$\theta(\tau, \omega_{EP}).$$

The cdf for the valuation $\theta(\tau, \omega_{EP})$, conditional on the type and the information publicly available in real time, is

$$\Pr\{\theta(\tau, \omega_{EP}) \leq \theta | \omega_{RT}\} \equiv F(\theta | \tau, \omega_{RT}).$$

We denote the conditional pdf $\partial / \partial \theta F(\theta | \tau, \omega_{RT})$ as $f(\theta | \tau, \omega_{RT})$.

The set of distributions over ex post valuations obeys a monotonicity and smoothness condition:

**Assumption 1 (First-order stochastic dominance—FOSD).** The distribution over ex post valuations is ordered by, and differentiable with respect to, ex ante type:

1. $F(\theta | \tau, \omega_{RT}) < F(\theta | \tau', \omega_{RT}) \quad \forall \tau > \tau', \forall \theta, \omega_{RT};$

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13This could easily be generalized, but assuming a very high density over a short interval of types should be a reasonable approximation to a point mass on a single type, and this setup allows us to distinguish increments anonymously—i.e. only by type—which is convenient.
2. \( \partial F(\theta | \tau, \omega_{RT}) / \partial \tau < 0 \); and

3. \( \exists M \in \mathbb{R}_+ \) s.t. \( |\partial F(\theta | \tau, \omega_{RT}) / \partial \tau| < M \) for a.e. \((\omega, \tau) \in \Omega_{RT} \times [\tau, \bar{\tau}]\) under the product measure.

In fact, this condition can be weakened so that there is a different FOSD ordering of the type space for each real time outcome, but the corresponding informational requirements for the aggregator may seem unrealistically onerous.

Our second assumption is a restriction on the set of DR curtailment policies:

**Assumption 2 (Merit Order Curtailment Policy).** We restrict attention to DR curtailment policies with a “Merit Order,” or cutoff, form: \( k(\tau, \omega_{RT}) = 1_{\{\tau \leq \hat{\tau}(\omega_{RT})\}} \). That is, in each real-time RT outcome \( \omega_{RT} \), the aggregator sends the curtailment signal to all increments with ex ante type less than some event-specific cutoff type, \( \tau = \hat{\tau}(\omega_{RT}) \).

We explain in Section 2.3 below how the curtailment policy can take the increments’ types as an argument, despite that the types are the increments’ private information.) The choice of DR dispatch policy \( \hat{\tau}(\cdot) \) determines the quantity of demand response, which is a random variable, whose value is realized ex post: \( DR(\omega_{EP}; \hat{\tau}(\omega_{RT})) \) (see Definition 3.1).

The combination of Assumptions 1 and 2 ensures that the type space \([\tau, \bar{\tau}]\) is “one-dimensional” in a key economic sense.\(^{14}\) Consider any pair of types \( \tau_2 > \tau_1 \) and any merit order allocation rule \( k(\tau, \omega_{RT}) = 1_{\{\tau \leq \hat{\tau}(\omega_{RT})\}} \). The marginal utility for type \( \tau_2 \) being switched from allocation \( k(\tau_1, \cdot) \) to allocation \( k(\tau_2, \cdot) \) is greater than the marginal utility for \( \tau_1 \) undergoing the same switch. So higher types value “higher allocations” (that is, being curtailed less) more than do lower types, a fact which allows the aggregator to “separate” the types. This essentially reduces DR contracting problem to a textbook single-stage screening problem by separability across events \( \omega_{RT} \in \Omega_{RT} \).

In a standard sequential screening problem (Courty and Li, 2000), a single agent’s ex post valuation is realized conditional on its ex ante type. Our problem is putatively dynamic, but because we exogenously specify how the increment makes the final consumption decision, it becomes effectively static, except that the increments’ first-stage utility function reflects the information dynamics.\(^ {15}\) In addition to the dynamic aspect, in our problem, our aggregator simultaneously derives benefit from a whole population of DR participants who have stochastic and possibly correlated valuations, rather than drawing a single agent from a common distribution. But, as reflects our application area, DR participants are “too small” to affect each other, so they are not strategically relevant to each other. These facts reduce our problem to a variation on a textbook screening problem, resulting in an expression for the optimal payment for a given curtailment policy, in Proposition 2.10 (“revenue equivalence”).

\(^{14}\)See Börgers (2010), chapter 5.6.

\(^{15}\)This “staticness” holds in a more specific sense than the more general result that sequential mechanisms can be reduced to a particular kind of static mechanism (Krähmer and Strausz, 2015). This reflects our concrete interest in demand response as embedded in our end-to-end problem.
This makes it straightforward to embed the contracting decision in the wholesale market problem.

For technical reasons, we also assume that the distribution $F(\theta|\tau, \omega_{RT})$ has constant support for all $\tau, \omega_{RT}$.

Finally, we make a very plausible simplifying assumption, that the highest type, $\overline{\tau}$, is “very high”:\footnote{We just mean that there are some units of demand whose valuation for power is high enough that they would not accept a reasonable payment curtailment in any contingency. For example, hospitals with life support units may have some quantity level at which their demand is, for all intents and purposes, perfectly inelastic.}

**Assumption 3.** For any optimal curtailment policy $\hat{\tau}^*(\cdot)$, $\overline{\tau} > \hat{\tau}^*(\omega_{RT})$.

\forall \omega_{RT} \in \Omega_{RT}.

### 2.2 Direct mechanisms

We assume that ex ante, DR increments are anonymous: the only distinguishing feature of each one is its type. Therefore, the result of contacting will be that, one way or another, each capacity increment is assigned a curtailment status, contingent only on its type, and the real-time dispatch outcome $\omega_{RT}$; this curtailment status function is its “allocation,” which we denote as $k(\tau, \omega_{RT})$. The contracting outcome will also involve a payment from the aggregator to the increment. Because both the aggregator and the increments are risk-neutral and we assume that the aggregator is capable of commitment, it makes no difference whether the payment is made ex post, or whether the same payment is made ex ante as an expectation over the ex post payments.

There might be many mechanisms by which such a contracting outcome could come about, but the Revelation Principle shows that any equilibrium contracting outcome can be achieved by a “direct revelation mechanism.”\footnote{Here we closely follow the development of Börgers (2010), chapter 2.}

**Definition 2.2.** A direct mechanism in our setting is a pair of functions

\[
([\bar{\tau}, \bar{\tau}] \ni \tau \mapsto k(\tau, \cdot) \in [0, 1])^{\bar{\omega}_{RT}}
\]

(where $[0, 1]^{\bar{\omega}_{RT}}$ is the set of functions having domain $\Omega_{RT}$ and codomain $[0, 1]$) and

\[
([\bar{\tau}, \bar{\tau}] \ni \tau \mapsto t(\tau) \in \mathbb{R})
\]

Here $k(\bar{\tau}, \cdot)$ is the real-time dispatch curtailment function, assigned ex ante to an increment reporting type $\bar{\tau}$, and $t(\bar{\tau})$ is the ex ante or expected gross payment made to an increment reporting $\bar{\tau}$.

The interpretation is that the increment of type $\tau$ gives a “report” of its type, $\bar{\tau} \in [\bar{\tau}, \bar{\tau}]$, and the aggregator commits to effect the corresponding allocation function and payment, $\langle k(\bar{\tau}, \cdot), t(\bar{\tau}) \rangle$. We call $t(\bar{\tau})$ the “gross ex ante payment” to an increment reporting type $\bar{\tau}$. It is “gross” because the aggregator may also collect a penalty from the increment ex post if it violates the curtailment signal, and we will net the latter quantity out of the aggregate payment.
2.2.1 The increments’ utility function in a direct mechanism

Here we display the increments’ utility function in a direct mechanism in a manner that reflects our demand response setting. In the following subsection, we show that any contracting equilibrium that allocates a curtailment function and an expected payment can be implemented by a direct mechanism—in particular, a “direct revelation mechanism.”

A capacity increment of infinitesimal magnitude \( d x \), with ex post valuation \( \theta \), pays the variable charge \( R d x \) or \( H d x \) (depending on its curtailment status) and enjoys utility \( \theta d x \) from consuming the quantity \( d x \) MWh. Henceforth, we normalize the quantities produces, pays or enjoys—DR, tariffs or expected utility—by dividing them by \( d x \); resulting in the corresponding quantity “per unit mass” of increments.

A priori, in its “outside option”, each increment has a retail service contract that permits it to consume if it pays the retail rate \( R \) $/MWh. This service contract is an option, in the financial sense, for physical delivery of a commodity at the point of service. Before the institution of a DR policy, this option only has value to its owner, since the commodity cannot be transferred and thus enjoyed by anyone else. But a DR policy establishes (in our model), by administrative fiat, that if the aggregator prevents an increment from exercising its service option when it otherwise would have, then the resulting reduction in consumption is treated as production of that same quantity by the aggregator.

An increment’s expected utility from holding its original service option is its “option value.” This is its consumption utility net of the retail price, provided that it is positive. Through contracting, the aggregator purchases the right to send a curtailment to each increment in every event where the reported type \( \tilde{\tau} \leq \tilde{\tau}(\omega_{RT}) \). The curtailment signal penalizes consumption by raising the effective price (exercise price) of the increment’s service option from the retail rate, \( R \), to an exogenously determined penalty rate, \( H > R \). This reduces both the option value and the quantity of consumption.

Generally, we assume that the aggregator collects the difference or “penalty fee” \( H - R \) if the increment consumes despite receiving the signal, but we are particularly interested in the special case of of direct load control, which is modeled by \( H = \infty \). In that case a penalty fee is never collected, because the increment always complies with the curtailment signal.

We quantify the increment’s utility when under contract as net of the outside option value:

**Definition 2.3** (Net option value from curtailment). Denote an increment’s change in ex post value given curtailment as

\[
L(\theta) \triangleq (\theta - H)^+ - (\theta - R)^+ \leq 0 . \tag{2}
\]

The lost option value for type \( \tau \), from perspective of real time event \( \omega_{RT} \) conditional on curtailment, is

\[
z(\tau, \omega_{RT}) \triangleq \int_{\Theta} L(\theta) f(\theta|\tau, \omega_{RT}) d\theta \leq 0 . \tag{3}
\]

The net option value of an increment of type \( \tau \), reporting type \( \tilde{\tau} \), given
curtailment policy \( k(\cdot, \cdot) \), is
\[
U(\tilde{\tau} | \tau) \triangleq \int_{\Omega_{RT}} z(\tau, \omega) k(\tilde{\tau}, \omega) dP_{RT}(\omega).
\]
(4)

We assume that increments have quasi-linear utility, so that it affords the following decomposition:

**Definition 2.4** (Direct mechanism representation of the increments’ expected utility function). An increment that reports its type as \( \tilde{\tau} \), given that it has true type \( \tau \), enjoys expected net utility
\[
u(\tilde{\tau} | \tau) \triangleq U(\tilde{\tau} | \tau) + t(\tilde{\tau}).
\]
(5)

### 2.3 The Revelation Principle

The Revelation Principle establishes that without loss of generality, we can restrict attention to “direct revelation mechanisms,” in which the increment reports its true type to the aggregator’s direct mechanism, and in which the incentive compatibility constraint is satisfied.

We quote Proposition 2.1 of Börgers (2010) with minor substitutions, the proof of which can be found there:

**Proposition 2.5** (Revelation Principle). For every mechanism \( \Gamma \) and every optimal increment strategy \( \sigma \) in \( \Gamma \), there is a direct mechanism \( \Gamma' \) and an optimal buyer strategy \( \sigma' \) in \( \Gamma' \) such that

1. The strategy \( \sigma' \) satisfies:
\[
\sigma'(\tau) = \tau \text{ for every } \tau \in [\tau, \overline{\tau}],
\]
i.e. \( \sigma' \) prescribes telling the truth;

2. For every \( \tau \in [\tau, \overline{\tau}] \), the curtailment allocation \( k(\tau, \cdot) \) and the payment \( t(\tau) \) equal the allocation function and the expected payment that result in \( \Gamma \) if the buyer plays her optimal strategy \( \sigma \).

**Definition 2.6** (Incentive compatibility). A direct mechanism in our problem is incentive compatible if truth-telling is optimal for every type; that is if:
\[
\tau \in \arg \max_{\tilde{\tau} \in [\tau, \overline{\tau}]} \{ U(\tilde{\tau} | \tau) + t(\tilde{\tau}) \}.
\]
(IC)

The previous result and definition allow us to simplify our problem, by restricting attention to direct revelation mechanisms which satisfy the (IC) constraint.\(^{18}\) An increment participating in an incentive compatible mechanism derives its equilibrium utility:

**Definition 2.7** (Equilibrium utility). For a given direct revelation mechanism, an increment of type \( \tau \) enjoys net expected utility
\[
u(\tau) \triangleq u(\tau | \tau) = U(\tau | \tau) + t(\tau)
\]
in the contracting equilibrium, i.e., when it truthfully reports its type in a direct mechanism.

\(^{18}\)We should note, however, that the Revelation Principle assumes away problems of multiple equilibria. However, this is not a problem for us, because we have already restricted the curtailment policy set in Assumption 2.
Another crucial constraint results from our assumption that the increments participate in the contracting scheme voluntarily:

**Assumption 4 (Individual rationality).**

\[ u(\tau) \geq 0 \quad \forall \tau \in [\underline{\tau}, \overline{\tau}]. \quad \text{(IR)} \]

This is to say that contracting cannot leave the increment worse off than in its outside option, which we normalize to zero. (Remember that \( U(\tilde{\tau}|\tau) \) is the change in consumption utility, and penalties, resulting from contracting.) Further, we assume that all increments participate, without loss of generality: for suppose an increment found it better not to participate, so that it is never curtailed and generates no DR for the aggregator, and receives no payment from the aggregator. This could be equivalently represented by \( k(\tau, \cdot) \equiv 0 \), and \( t(\tau) = 0 \). In our model, this allocation makes it so the increment makes no contribution to the aggregator’s objective, and by definition of the outside option, the increment’s net utility is also zero. So non-participation can be modeled as the degenerate form of participation just mentioned, and without loss of generality we can enforce the assumption that all increments participate and achieve nonnegative net utility.

### 2.4 The optimal payment \( T \) needed to effect a demand response policy

We are now ready to derive an expression for the aggregate payment \( T \) as a function of the curtailment policy.

**Lemma 2.8** (Continuity and differentiability of equilibrium consumption utility with respect to true type). The increment’s utility function, as a function of its true type while holding its report constant, is differentiable, Lipschitz continuous, and thus absolutely continuous, for any given report \( \tilde{\tau} \). Therefore there exists an integrable function \( b(\tau) \) such that \[ \frac{1}{\tilde{\tau}} U(\tilde{\tau}|\tau) \leq b(\tau), \text{ and } U(\tilde{\tau}|\cdot) \text{ is uniformly continuous, } \forall \tilde{\tau}. \]

**Proof.** See appendix. \( \square \)

At this point it would be standard to state a necessary condition that incentive compatibility places on the curtailment allocation \( k(\cdot, \cdot) \). However, since we are restricting attention to the class of cutoff policies \( \{\tilde{\tau}(\cdot)\} \), we will skip this, and show in Proposition 2.12 that all cutoff policies can be implemented, given the appropriate payment.

**Lemma 2.9.** Incentive compatibility implies that the equilibrium net expected utility \( u(\tau) \) is decreasing in ex ante type, that \( u(\tau) \) is absolutely continuous,
and that
\[
    u(\tau) = u(\tau) - \int_{\tau}^{\hat{\tau}} \frac{\partial}{\partial s} U(\hat{\tau}|s) \bigg|_{s=\tau} ds
\]
\[= u(\tau) + \int_{\tau}^{\hat{\tau}} \int_{\Omega} \int_{\Theta} L' (\theta) \frac{\partial F(\theta|\tau, \omega)}{\partial s} k(s, \omega) dP(\omega) ds
\]
\[= u(\tau) - \int_{\tau}^{\hat{\tau}} \int_{\Omega} \int_{\Theta} L(\theta) \frac{\partial f(\theta|\tau, \omega)}{\partial s} k(s, \omega) dP(\omega) ds.
\]

**Proof.** The claim follows from Milgrom and Segal (2002), Theorem 2, the conditions of which we have established in Lemma 2.8. Since the integrand in line (7) is nonnegative, and \(\tau\) is the lower limit of integration, \(u(\tau)\) is nonincreasing.

**Proposition 2.10** (Necessary condition on the payment for Incentive Compatibility / Revenue Equivalence). Under a merit-order curtailment policy, incentive compatibility requires that the ex ante gross payment to an increment of type \(\tau\) be
\[
t(\tau) = u(\tau) - \int_{\Omega_{RT}} z(\hat{\tau}(\omega), \omega) k(\tau, \omega) dP_{RT}(\omega).
\]

**Proof.** See appendix.

The next proposition shows that when the aggregator optimizes its profit, \(u(\tau)\) is zero. So Equation 9 can be interpreted as saying that, for each real time state, every curtailed increment is paid the reservation utility of the highest curtailed increment in that state.

**Proposition 2.11** (Individual Rationality “binds at the top”). Given incentive compatibility, individual rationality holds only if \(u(\tau) \geq 0\). Maximization of the aggregator’s profit implies that \(u(\tau) = 0\).

**Proof.** Proposition 2.9 implies that \(u(\tau)\) is nonincreasing. Since the aggregator’s payment to each increment includes the constant term \(u(\tau) \geq 0\), the aggregator maximizes profit by reducing \(t(\tau)\) until \(u(\tau) = 0\).

**Proposition 2.12** (Sufficient Condition for Incentive Compatibility). Any merit-order curtailment policy is rendered incentive compatible, when the corresponding payment is as in Proposition 2.10.

**Proof.** See appendix.

This proposition shows that the DR contracting problem is separable over \(\Omega_{RT}\).
2.4.1 The payment form suggests a real time implementation of DR contracting

The formula in Proposition 2.10 for the payment (setting \( u(\tau) = 0 \)), has an intuitive interpretation, corresponding to what we call the “real time implementation” of \( k(\cdot, \cdot) \). Proposition 2.12 shows that under First Order Stochastic Dominance (Assumption 1), merit order curtailment guarantees the satisfaction of any linking constraints across random events \( \omega_{RT} \) that are implied by incentive compatibility, so that those linking constraints can be discarded. The aggregator’s DR contracting problem is therefore separable across RT states \( \omega_{RT} \), and we can think of the aggregator as making a separate DR purchase in each \( \omega_{RT} \). In the merit order setting the curtailment policy for each fixed real time outcome is binary and decreasing over the type argument: i.e., it is constant over all curtailed types. This is because incentive compatibility forces the aggregator to pay each curtailed increment the reservation value of the highest curtailed type, since any increment could impersonate that type. In the mechanism we focus on, the aggregator and the increment make ex ante contracting commitments. However, these considerations show that the aggregator can achieve the same distribution over curtailment outcomes without ex ante contracting, by waiting until the realization of \( \omega_{RT} \) and quoting a curtailment price offered for curtailment in that state, equal to the lost expected utility of the marginally curtailed increment. All types accepting curtailment in event \( \omega_{RT} \) would receive that same payment.\(^{19}\)

2.5 Expressing the net payment as a linear functional of the curtailment policy

In order to solve the end-to-end problem, we need expressions of a certain form for the aggregate payment, i.e. integrated over the DR population.

**Definition 2.13 (Payment).** The aggregate net payment is the integral of the type-specific payment over the increment population, net of penalty receipts:

\[
T \triangleq \int_{\tau}^{\tau} t(\tau)g(\tau) d\tau - (H - R) \int_{\Omega_{RT}} \int_{\tau}^{\tau} \mathbb{1}\{\tau \leq \hat{\tau}(\omega_{RT})\} \mathbb{1}\{\theta(\tau, \omega_{EP}) > H\} g(\tau) d\tau .
\]

(aggregate expected penalty fee receipts)

**Proposition 2.14 (Expressing the aggregate payment as an expectation of a linear functional of the increment population).** Defining the virtual net payment to type \( \tau \) in event \( \omega_{RT} \) as

\[
\psi(\tau, \omega_{RT}) \triangleq - \left( z(\tau, \omega_{RT}) + \frac{G(\tau)}{g(\tau)} \frac{\partial}{\partial \tau} z(\tau, \omega_{RT}) \right) - (H - R) \Pr\{\theta(\tau, \omega_{EP}) \geq H|\omega_{RT}\},
\]

\(^{19}\)This dispatch policy structure is less restrictive than it may appear to be, because we can augment the RT state with a non-payoff-relevant randomization variable. However, analysis of the first order conditions shows that the aggregator does not benefit from randomization.

\(^{20}\)See the discussion at the end of Börgers (2010), chapter 2.2.
The aggregate payment can be expressed as

\[ T = \int_{\Omega} \int_{\hat{\tau}} \psi(\tau, \omega_{RT}) g(\tau) d\tau \, dP(\omega_{RT}). \]  

(10)

**Proof.** See appendix.  

We call \( \psi(\tau, \omega_{RT}) \) the “virtual net payment” to type \( \tau \) for curtailment in event \( \omega_{RT} \). In addition to penalty fee receipts, the marginal aggregate payment from raising the marginally curtailed type to \( \tau \) in event \( \omega_{RT} \) has two components: the aggregator must pay the marginal increments their own reservation price \( \varepsilon(\tau, \omega_{RT}) g(\tau) \); and it must also raise the payment it makes to infra-marginal types to that same reservation price, incurring a payment \( G(\tau) \frac{\partial}{\partial \tau} \varepsilon(\tau, \omega_{RT}) \), because the infra-marginal types can impersonate the marginal type. The same economic insight arises when we interpret the product rule in the first order condition for the elementary monopoly pricing problem. The expression \( \psi(\tau, \omega_{RT}) \) attributes both of these components to the marginal type, so that we obtain the marginal change in the aggregate payment from recruiting type \( \tau \). This helps us write the first order conditions for the end-to-end problem.

### 2.6 Monopsony under-purchase of DR

Examining the formula for \( \psi(\tau, \omega_{RT}) \) in Proposition 2.14, we note that since first order stochastic dominance guarantees that \( \frac{\partial}{\partial \tau} \varepsilon(\tau, \omega) \leq 0 \), the virtual payment is greater than the increment’s lost utility from curtailment. This implies that the aggregator will purchase less demand response than the efficient level, because at the efficient level, the cost of DR purchase from the marginal increment must be equal to that increment’s lost utility. The efficient level can be achieved either through perfect competition among aggregators, or if the aggregator had perfect information of increment types; although in the latter case, the DR participants would be worse off. We discuss this further in the conclusion.

### 3 Analyzing the end-to-end problem

In the previous section, we derived useful expressions for the cost of recruiting DR. To maximize its profit, the aggregator balances costs against benefits. Having determined the cost of DR as a function of the policy, we now analyze the aggregator’s problem as a two-stage problem: first, we characterize the optimal dispatch of DR, conditional on an arbitrary DA offer policy \( q \). Then, holding the DR policy at its optimal setting as a function of \( q \), we optimize \( q \).

---

21 It turns out that this formula holds for general \( k(\cdot, \cdot) \) satisfying incentive compatibility (i.e. we would integrate the above expression over the whole population, multiplying the integrand by \( k \)), not just for the cutoff form \( \hat{\tau} \), but this is not an immediate concern of ours here, so we leave this result unproved.
3.1 Benefit from DR dispatch

Recall the informal sketch of the objective from Equation (1):

\[
\max_{q,DR,T} J(q,DR,T) = \max_{q,DR,T} \mathbb{E} [pq + a(\text{DR} + s - q)^+ - b(q - \text{DR} - s)^+] - T
\]

We now give the definition of the DR production quantity:

**Definition 3.1** (DR quantity, \(DR(\hat{\tau}(\cdot), \omega_{EP})\)). The quantity of DR in ex post event \(\omega_{EP}\) is the measure under \(g\) of increments that forego consumption as a result of receiving the curtailment signal:

\[
DR(\hat{\tau}(\omega_{RT}); \omega_{EP}) \Delta \int_{\omega_{RT}}^{\hat{\tau}(\omega_{RT})} \mathbb{I} \{\tau : R \leq \theta(\tau, \omega_{EP}) \leq H\} g(\tau) d\tau
\]

We assume that the valuation process allows the application of Fubini’s theorem to the DR process:

**Assumption 5** (Fubini property of ex post DR process). We assume that for all events \(\omega_{RT} \in \Omega_{RT}\), and all subsets \(B \subset [\tau, \tau]\)

\[
\mathbb{E}_{EP} \left[ \int_{B} \mathbb{I} \{\theta(\tau, \omega_{EP}) \in [R,H]\} g(\tau) d\tau \bigg| \omega_{RT} \right] = \int_{B} \mathbb{P}_{EP}(\theta(\tau, \omega_{EP}) \in [R,H] | \omega_{RT}) g(\tau) d\tau
\]

\[
= \int_{B} y(\tau, \omega_{RT}) g(\tau) d\tau
\]

Here we define \(y(\tau, \omega_{RT})\) as the “expected DR yield” of type \(\tau\) conditional on event \(\omega_{RT}\). Also, recall that the DA offer policy is a function of the day-ahead information: \(q : \Omega_{DA} \rightarrow \mathbb{R}\). We will analyze the aggregator’s end-to-end problem in two special cases.

3.2 Example 1: purely idiosyncratic valuation shocks

In this section, we assume that conditional on the real-time outcome, uncertainty regarding the ex post valuation process is i.i.d. noise.

**Assumption 6.** The valuation process decomposes as

\[
\theta(\tau, \omega_{EP}) = m(\tau, \omega_{RT}) + e(\tau, \omega_{EP})
\]

where \(m(\tau, \omega_{RT})\) is a deterministic function, \(e(\tau, \omega_{EP})\) is i.i.d. conditional on \(\omega_{RT}\) over \(\tau\), and \(\mathbb{E}[e(\tau, \omega_{EP})] = \mathbb{E}[e(\tau, \omega_{EP}) | \omega_{RT}] = 0, \forall \omega_{RT}, \omega_{EP}\). Each \(e(\tau, \cdot)\) has conditional cdf \(\Phi(e; \omega_{RT})\), and conditional pdf \(\varphi(e; \omega_{RT})\).
Remark 3.2. Assumptions 1 (FOSD) and 6 (idiosyncratic noise) jointly imply that \( m(\cdot, \omega_{RT}) \) is increasing, for each \( \omega_{RT} \). The common conditional cdf of valuation shocks can be written as \( \Pr(\theta \leq \tau | \omega_{RT}) = \Phi(z - m(\tau, \omega)|\omega_{RT}) \).

This determines the virtual payment function \( \psi(\tau, \omega_{RT}), \) the formula for which we omit. Further, our assumption of constant support for \( \{ \theta_{\tau} \} \) implies that each \( \theta_{\tau} \) has full support on \( \mathbb{R} \) for every \( \omega_{RT} \).

We make the following assertion without proof:

Remark 3.3. Under Assumption 6, the DR output is almost surely deterministic, conditional on \( \tau(\omega_{RT}) \):

\[
DR(\hat{\tau}(\omega_{RT}); \omega_{EP}) = \int_{\tau(\omega_{RT})}^{\xi(\omega_{RT})} \mathbb{I}\{\theta(\tau, \omega_{EP}) \in [R, H]\} g(\tau) d\tau = \mathbb{E}_{\omega_{EP}} \left[ \int_{\tau(\omega_{RT})}^{\xi(\omega_{RT})} \mathbb{I}\{\theta(\tau, \omega_{EP}) \in [R, H]\} g(\tau) d\tau \right] = \int_{\tau(\omega_{RT})}^{\xi(\omega_{RT})} \Pr_{\omega_{EP}}(\theta(\tau, \omega_{EP}) \in [R, H]|\omega_{RT}) g(\tau) d\tau = \int_{\tau(\omega_{RT})}^{\xi(\omega_{RT})} y(\tau, \omega_{RT}) g(\tau) d\tau.
\]

See Al-Najjar (1995) for further discussion on this issue. Our purpose in making Assumption 6 was simply to license the above result; the reader may take this consequence as the operative assumption instead.

Having derived expressions for the DR production quantity and the payment contribution in each RT outcome \( \omega_{RT} \), we can now plug them into the stylized objective of Equation (1). Note that the real-time contribution to the ex ante component is the integrand from Equation (10).

Definition 3.4 (Aggregator’s objective, and its constituent parts). The real-time contribution to the aggregator’s ex ante objective in event \( \omega_{RT} \) is

\[
J_{RT}(q, \hat{\tau}(\omega_{RT}); \omega_{RT}) = p q + a \left( \int_{\tau(\omega_{RT})}^{\xi(\omega_{RT})} y(\tau, \omega_{RT}) g(\tau) d\tau + s - q \right)^{+} - b \left( q - \int_{\tau(\omega_{RT})}^{\xi(\omega_{RT})} y(\tau, \omega_{RT}) g(\tau) d\tau - s \right)^{+} - \int_{\tau(\omega_{RT})}^{\xi(\omega_{RT})} \psi(\tau, \omega_{RT}) g(\tau) d\tau; \tag{11}
\]

The day-ahead contribution is the expectation over the real time contribution:

\[
J_{DA}(q, \hat{\tau}(\cdot); \omega_{DA}) \doteq \mathbb{E}_{\omega_{RT} \mid \omega_{DA}} [J_{RT}(q, \hat{\tau}(\omega_{RT}); \omega_{RT}) | \omega_{DA}]; \tag{12}
\]

And the aggregator’s objective, full stop, is the expectation over the DA contribution:

\[
J_{EA}(q(\cdot), \hat{\tau}(\cdot)) \doteq \mathbb{E}_{\omega_{DA}} [J_{DA}(q(\omega_{DA}), \hat{\tau}(\cdot); \omega_{DA})]. \tag{13}
\]
Approaching the problem as a two-stage decision problem with recourse, we first derive a first order necessary condition for optimizing Equation (11) with respect to the optimal DR dispatch \( \hat{\tau}(\omega_{RT}) \) given a DA offer \( q \). We denote this optimal recourse policy as \( \hat{\tau}^*(\omega_{RT}|q) \). Since we will optimize \( q \) below, we also occasionally drop the "\(|q\)" or "conditioned on \( q \)" argument for brevity.

### 3.2.1 First order necessary conditions for \( \hat{\tau}^*(\omega_{RT}|q) \) in example 1

To make the first order conditions and consequent results easier to read, we define the following quantities:

**Definition 3.5.** Under Assumption 6 (idiosyncratic noise), we define the marginal type that must be curtailed to cancel out a nonnegative RT imbalance, given wind realization \( s \), as

\[
DR^-1(q-s; \omega_{RT}) = \begin{cases} 
-\infty & \text{if } q < s \\
\tau & s.t. DR(\tau, \omega_{RT}) = q - s, \text{if } DR(\tau, \omega_{RT}) \geq q - s \geq 0 \\
\infty & \text{if } q - s \geq DR(\tau, \omega_{RT}).
\end{cases}
\]

**Definition 3.6.** The marginal cost of DR from type \( \tau \) (i.e. the dollar amount the aggregator must pay to curtail increments of type \( \tau \), per unit of resulting DR yield), is  

\[
c(\tau; \omega_{RT}) = \frac{\psi(\tau, \omega_{RT})}{y(\tau, \omega_{RT})}.
\]

We will occasionally suppress \( \omega_{RT} \) arguments for \( DR, c, \) etc. for conciseness.

**Proposition 3.7** (Optimal DR curtailment policy). The first order necessary condition requires that if \( \hat{\tau}^*(\omega_{RT}) > \tau \), the following condition holds (here \( \partial(\cdot) \) denotes the subgradient mapping):

\[
\psi(\hat{\tau}^*)g(\hat{\tau}^*) \in \partial \left( \int \left[ y(\tau; \omega_{RT})1_{\tau \leq \cdot} g(\tau) d\tau + s - q \right]^+ 
- b(q - s - \int \left[ y(\tau; \omega_{RT})1_{\tau \leq \cdot} g(\tau) d\tau \right]^+) \right)
\]

which implies

\[
c(\hat{\tau}^*) \in \{a\}1_{DR(\hat{\tau}^*+\tau)>q} + \{b\}1_{DR(\hat{\tau}^*)+\tau<q} + [a, b]1_{DR(\hat{\tau}^*)+\tau=q}.
\]

If \( c(\cdot) \) is monotonically increasing, this requires that (suppressing the \( \omega_{RT} \) argument of \( DR^-1 \))

\[
c(\hat{\tau}^*(\omega_{RT})) = \begin{cases} 
c(DR^-1(q-s)) & \text{if } c(DR^-1(q-s)) \in (a, b), \\
a & \text{if } c(DR^-1(q-s)) < a, \text{ and} \\
b & \text{if } c(DR^-1(q-s)) > b.
\end{cases}
\]

**Proof.** Subdifferentiation of the previous formula for the aggregator's objective with respect to \( \hat{\tau}^* \). \( \square \)
Assuming that \( c \), the marginal cost of DR, is monotonic, in the optimal DR recourse policy \( \hat{\tau}^* (\omega_{RT} | q) \), the aggregator exactly meets its commitment if the marginal cost of DR needed to do so is strictly between the two imbalance prices, \( a \) and \( b \).\(^{22}\) Otherwise, if the marginal cost of DR needed to do so would be greater than \( b \), the aggregator curtails all increments up to the “upper critical valuation,” \( c^{-1}(b) \). If the marginal cost of curtailing up to the point of zero imbalance is below \( a \), then the aggregator produces more than it offered day ahead, curtailing all increments up to the “lower critical type,” \( c^{-1}(a) \). Increments with type below the minimum lower critical type, \( \tau < \inf_{\omega_{RT}} c^{-1}(a(\omega_{RT}); \omega_{RT}) \) are curtailed in all wind outcomes; increments with valuation above the maximum upper critical type are not curtailed in any wind outcome. These features of the real time curtailment decision are depicted in the upper left panel of Figure 2, although that figure is further specialized to Example 2, introduced below. That figure is the monopsony version of the standard presentation of optimal monopoly pricing, but with a piecewise constant benefit function, reflecting our aggregator’s newsvendor-style revenue function.

Decomposing events by imbalance status, the first order conditions for monotonic \( c \) admit the following formulation,\(^{23}\) which is also helpful for the optimization of \( q \).

\[
\hat{\tau}^* (\omega_{RT} | q) = \begin{cases} 
DR^{-1}(q-s) & \text{if } DR(\hat{\tau}^* (\omega_{RT}|q)) = q-s, \\
\hat{c}^{-1}(b) & \text{if } DR(\hat{\tau}^* (\omega_{RT}|q)) < q-s, \text{ and} \\
\hat{c}^{-1}(a) & \text{if } DR(\hat{\tau}^* (\omega_{RT}|q)) > q-s.
\end{cases}
\] (15)

3.2.2 The optimal day-ahead offer policy, \( q^* \)

To find the optimal \( q^* \), we first plug the expression for the optimal DR dispatch in Equation (15) into the aggregator’s objective, conditional on \( \omega_{DA} \):

\[
J_{DA} (q, \hat{\tau}^* (|q); \omega_{DA}) = \sup_{\hat{\tau}(\cdot)} J_{DA} (q, \hat{\tau}(\cdot); \omega_{DA})
\]

\[
= pq + E_{\omega_{RT} | \omega_{DA}} \left[ a \left( \int_{\tau}^{\hat{\tau}^* (\omega_{RT} | q)} y(\tau) g(\tau) \, d\tau + s - q \right) \right] \\
- E_{\omega_{RT} | \omega_{DA}} \left[ b(q-s - \int_{\tau}^{\hat{\tau}^* (\omega_{RT} | q)} y(\tau) g(\tau) \, d\tau) \right] \\
- E_{\omega_{RT} | \omega_{DA}} \int_{\tau}^{\hat{\tau}^* (\omega_{RT} | q)} \psi(\tau; \omega_{RT}) g(\tau) \, d\tau .
\] (16)

We obtain the first order condition by differentiation:

\(^{22}\)We will argue in a followup work that the cost \( c \) should typically be monotonic. Another setting perhaps worth considering is for \( c \) “single-troughed,” i.e. quasi-convex, so that \(-c\) is single-peaked. This might obtain if it is prohibitively expensive to obtain DR from types low valuation types, because their yield is too low. We may consider the optimality conditions for this case elsewhere.

\(^{23}\)If \( c \) is weakly increasing, then the equality is replaced with \( c \), and \( c^{-1} \) is interpreted as the set-valued preimage function.
Proposition 3.8 (First order necessary condition for optimal day-ahead offer). The optimal day-ahead offer, \( q^* \), conditional on \( \omega_{DA} \), satisfies the following condition:

\[
p = \mathbb{E}[a|DR(\hat{\tau}^* + s > q^*)] \Pr(\hat{\tau}^* + s > q^*)
+ \mathbb{E}[b|DR(\hat{\tau}^* + s < q^*)] \Pr(\hat{\tau}^* + s < q^*)
+ \mathbb{E}[c(\hat{\tau}^*)|DR(\hat{\tau}^* + s = q)] \Pr(\hat{\tau}^* + s = q^*) .
\]

Proof. See appendix.

This is not an explicit solution; we present this condition because of its clear economic interpretation: The aggregator increases its DA offer quantity until the marginal change in expected real-time expenditures (imbalance prices plus payments to DR) rises to meet the marginal DA revenue. If the RT imbalance prices \( a \) and \( b \) are known day-ahead, then Figure 1 and the accompanying discussion in Section 3.3.1 apply to this case. That is, we show in that section how the aggregator’s problem can be viewed as an elaboration of the elementary monopsony pricing problem.

3.3 Example 2: parameterized uniform distributions

Now we consider a simple concrete example, where renewable power output and increment types are distributed uniformly, valuation noise is degenerate (i.e. nonexistent), the imbalance prices are known day-ahead, the lowest valuation is equal to the retail rate, and the aggregator employs direct load control. In this case, we can derive formulas for the aggregator’s optimal policy and its relevant features, in order to display the solution graphically, and obtain quantitative sensitivity results. The main objects in the model are

\[
g(\tau) = \frac{d}{d\tau} G(\tau) = N \mathbb{1}_{\tau \in [\tau, \tau]}
\]

\[
s \sim \text{Uniform}[0, \tau]
\]

\[
\theta_\tau \equiv \tau \text{ (degenerate distribution at } \tau) \tag{17}
\]

\[
R = \tau
g H = \infty .
\]

With no valuation noise, this model does not satisfy our Assumption 1 (FOSD). But we only needed Assumption 1 in order to prove Lemma 2.8 (monotonicity and differentiability of the equilibrium utility with respect to true type), which we can directly verify. In our current setting,

\[
U(\hat{\tau}|\tau) = -\int_{\omega_{RT}} \tau k(\hat{\tau}, \omega_{RT}) d\mathbb{P}(\omega_{RT}) = -\tau \Pr(\hat{\tau} \leq \hat{\tau}(\omega_{RT})), \tag{18}
\]

\[
\therefore \frac{\partial}{\partial \tau} U(\hat{\tau}|\tau) = -\Pr(\hat{\tau} \leq \hat{\tau}(\omega_{RT})) \leq 0 . \tag{19}
\]

The remaining lemmas and propositions of Section 2 therefore follow.

Before obtaining formulas for the payment, etc., we can simplify the model by reparameterization. First, by a change of variables on \( \tau \), we normalize
the valuations to express them as net of the retail rate, and set $\tau = R = 0$. This specification results in the following model quantities: $\forall \tau \in [\underline{\tau}, \bar{\tau}],$

\[
\begin{align*}
  z(\tau) &= \tau \\
  \frac{\partial}{\partial \tau} z(\tau) &= 1 \\
  G(\tau)/g(\tau) &= \tau \\
  \psi(\tau) &= 2\tau \\
  y(\tau) &= \psi(\tau) \\
  c(\tau) &= 2\tau .
\end{align*}
\]

Instead of expressing the problem in terms of $N$, the DR population size, and $\bar{\tau}$, the VER nameplate capacity, we write the problem in terms of the parameter $\nu = N/(\bar{\tau}(\bar{\tau} - \tau))$: the density of increment-valuations per dollar, per MW nameplate capacity. Now power (and energy) quantities are expressed as a fraction of the nameplate capacity: the DA offer quantity now takes the form $q \leftarrow q/\bar{s}$, the fraction of VER nameplate capacity bid day ahead; DR quantities are in the same units; the random variable $s$ is reparameterized as $s \leftarrow s/\bar{s}$, now the wind realization's cdf value. Correspondingly, the aggregator's profit is now denominated in dollars per unit VER nameplate capacity, per hour.

Following the same two-stage solution method as before, we first consider the optimization of the real-time demand response dispatch. In this setting,

\[
\text{DR}(\hat{\tau}) = \nu \hat{\tau} .
\]

The aggregator’s “real time objective” is:

\[
J_{RT}(q, \hat{\tau}(\omega_{RT}); \omega_{RT}) = p q + a(s + \nu \hat{\tau}(\omega_{RT}) - q) - \frac{b(\bar{s} - \nu \hat{\tau}(\omega_{RT}))}{\nu} - T(\omega_{RT}) ,
\]

where the last term the net ex post payment in event $\omega_{RT}$, i.e. $\int_{\underline{\tau}}^{\hat{\tau}(\omega_{RT})} \psi(\tau, \omega_{RT}) g(\tau) d\tau$ from Proposition 2.14:

\[
T(\omega_{RT}) = \nu \int_{0}^{\hat{\tau}(\omega_{RT})} 2\tau d\tau = \nu \hat{\tau}(\omega_{RT})^2.
\]

The first order conditions for $\hat{\tau}(\omega_{RT}|q)$ give us that in the optimal recourse curtailment policy, Equation (15) takes the form

\[
\hat{\tau}^*(\omega_{RT}|q) = \begin{cases} 
  a/2 & \text{if } (q - s)/\nu < a/2 \equiv s > q - va/2 \\
  (q - s)/\nu & \text{if } a/2 < (q - s)/\nu < b/2 \equiv q - vb/2 < s < q - va/2 \\
  b/2 & \text{if } (q - s)/\nu > b/2 \equiv s < q - vb/2 .
\end{cases}
\]

The derivation of model parameters from elasticity estimates must be done before the change of variables, because an elasticity is a ratio involving the retail price, which we are eliminating from the problem.
That is, the lower and upper critical valuations mentioned above are
\[ c^{-1}(a) = a/2 \] and \[ c^{-1}(b) = b/2 \] respectively. Referring back to Equation (14), which applies here as well, we see that the optimal marginal DR cost is set equal to the imbalance cost that obtains given the resulting DR quantity. A marginal analysis of this optimization is depicted in the upper left panel of Figure 1. Henceforth we assume that \( \hat{\tau} \) is held at its optimal recourse value given \( q \), and proceed to analyze the “first stage” problem, \( \max_q J_{DA}(q, \hat{\tau}(\cdot|q); \omega_{DA}) \).

For a given policy and \( \omega_{DA} = (p, a, b) \), there are three possible types of event with respect to wind: shortfall, zero imbalance, and overproduction. To calculate the expected value of the objective over the real-time information \( s \), we define the wind levels that constitute breakpoints between these regimes:

\[ d ≜ \min(1, \max(0, q - \nu b/2)) \quad (26) \]

\[ e ≜ \min(1, \max(0, q - \nu a/2)) \quad (27) \]

Since \( s \) is now normalized to refer to wind realization’s cdf value, \( d \) is the probability of shortfall, and the level \( e \) is one minus the probability of overproduction.\(^{25}\) Plugging in the optimal \( \hat{\tau}(\cdot|q) \) and taking the the expectation over \( s \), we get

\[
\max_q J_{DA}(q, \hat{\tau}(\cdot|q); \omega_{DA}) = \max_q p q - \int_0^d \text{payment} \frac{\nu(b/2)^2}{2} + \frac{b(q - \nu b/2 - s)}{2} ds \\
- \int_d^e \text{payment} \frac{\nu(q - s)^2}{2} ds - \int_e^1 \text{payment} \frac{\nu(a/2)^2}{2} - \frac{\nu(s + \nu a/2 - q)}{2} ds.
\]

This “first-stage” (DA-stage) objective is concave (it is a day ahead expectation with respect to \( s \) of a concave real time benefit function), and piecewise polynomial in \( q \), with breakpoints where \( d \) and \( e \) hit zero or one. The combinations of possible sets of events based on parameter values generate many cases.

**Proposition 3.9** (Solution to the end-to-end problem, Example 2). The solution to the aggregator’s end-to-end problem in Example 2 is presented in Table 1.

By “solution,” we mean that we exhibit the optimal policies \( q \) and \( \hat{\tau} \), as a function of \( \omega_{DA} \) and \( \omega_{RT} \) respectively. If the aggregator is to publish a menu of contract choices ex ante, it would do this by taking expectations of the curtailment status, and payment, over \( \Omega_{DA} \), according to average market statistics. We provide an example of this in Section 3.3.2.

\(^{25}\)For example, \( d \) is the probability that wind is less than \( q - \nu b/2 \), which would imply that the wind quantity plus the maximum economical amount of DR be less than the offer quantity \( q \).
3.3.1 Graphical marginal analysis of the optimization of DR and q in Example 2

We can obtain more intuition regarding the optimization of the aggregator’s policy by considering the graphical depiction of the marginal analysis of the aggregator’s problem, in Figure 1.\textsuperscript{26} With respect to DR, the aggregator is a monopsonist, i.e., the sole buyer in a commodity market with many sellers. In each real-time realization $\omega_{\text{RT}}$, it faces the same marginal purchase cost curve (the marginal cost of DR, Definition 3.6) and it makes the optimal real-time curtailment decision by ensuring that its marginal cost of DR is between the imbalance prices that it faces on the margin. The marginal benefit from purchasing the $x$th MW of DR is denoted by $mb(x)$, a piecewise constant decreasing function depicted in black. The marginal cost from purchasing the $x$th MW of DR is equal to the virtual payment to the marginally curtailed increment, which is increasing: $c(x) = 2x / \nu$. The optimal curtailment quantity is the point on the quantity axis where the marginal benefit curve crosses the marginal cost curve. Projecting the intersection of the two curves onto the $y$ axis, we obtain the optimal DR recourse cost associated with that real time outcome, which we depict with a red circle.

Next we step back to the DA optimization of $q$. At the DA stage, the aggregator foresees that a random RT outcome will realize, at which point it will take a DR recourse action in the manner just described. In the current example, the imbalance prices $a$ and $b$ are known day-ahead, and the only random variable at the DA stage is the wind, $s$. The distribution over wind outcomes (here with pdf $h$), together with the aggregator’s choice of $q$, induces a distribution over breakpoints in the marginal benefit curve,

\textsuperscript{26}In that figure, we slightly abuse notation: there, $c(x) \triangleq c(DR^{-1}(x))$, and $\tau(x) \triangleq DR^{-1}(x)$.
Figure 2: Supply curves.

\( h(q - s) \), which we depict under the \( x \) axis of the last three panels of Figure 1. Adjusting \( q \) slides the pdf of breakpoints along the \( x \) axis. We depict the distribution over marginal benefit curves induced by a choice of \( q \) as a regularly spaced finite sample from it (curves in light gray). Assuming optimal RT recourse, a choice of \( q \) induces a distribution over recourse costs, which we depict in red and pink on the \( y \) axis. This distribution has a density component corresponding to the virtual payment to the marginally curtailed type when there is no imbalance, as well as two point-masses, at \( a \) and \( b \), depicted as red circles with area proportionate to their probability. The aggregator's optimal DA offer, \( q^* \), sets the expected recourse cost equal to the DA revenue. That is, imagining that gravity is pulling the recourse cost distribution to the right, the optimal \( q^* \) balances that distribution on the DA price, \( p \). We depict the result of this process for three different samples of \((p, a, b)\) in the last three panels of Figure 2.

The result of optimizing \( q \) contingent on DA information can be represented as a supply curve. However, in Example 2, since we assume that the imbalance prices \( a \) and \( b \) are revealed simultaneously with \( p \), the policy mapping DA information to the DA offer quantity \( q \) is actually a "supply surface." We display several representative slices of this surface in Figure 2. In this figure, we consider three cases where the imbalance prices are equal to the day ahead price, plus and minus a premium \( \epsilon \). As \( \epsilon \) increases in this figure, the shortfall penalty and the overproduction rate are increased. Both of these effects move in the same direction, encouraging the producer to bid a smaller fraction of nameplate. (A higher overproduction payment encourages the producer to bid less, because it reduces the producer's opportunity cost in scenarios where its supply exceeds its bid.) In future research we will discuss how one can solve the end-to-end problem numerically, via simulation, in a general model. The output of such an optimization can be offered as a supply curve in an ISO auction market.
3.3.2 Graphical depiction of the contract menu

Figure 1 characterizes the aggregator’s optimal DA action, \( q^* \), and RT recourse, \( \hat{\tau}^* \). Stepping back to the ex ante stage, we consider the ex ante contract menu that would implement \( \hat{\tau}^*(\cdot) \). This requires some assumption of market statistics over \( \Omega_{DA} \). We consider the example:

\[
\begin{align*}
p &\sim \text{Uniform}[10, 100] \\
a &= (1 - \delta)p \\
b &= (1 + \delta)p \\
\delta &\sim \text{Uniform}[0.1, 0.9].
\end{align*}
\]

We also assume that \( \nu = 1/100 \), which is derived from the following parameter choices under linear demand:\footnote{Note that \( \nu = N/(\delta(\tau - \overline{\tau})) = g(R)/\delta \). The elasticity at the retail rate is \( \eta(R) = g(R)/R \). The parameter chosen yield \( \nu = 1/100 \).}

\[
\begin{align*}
\bar{s} &= 100 \text{ MW} \\
R &= 30/\text{MW} \quad \text{(generation component of the retail price)} \\
N &= D(R) = 100 \text{ MW} \quad \text{(aggregate demand at } R = 30) \\
\eta(R) &= 0.3 \quad \text{(elasticity at } R = 30). 
\end{align*}
\]

In Figure 3, we plot the probability of curtailment, as well as the ex ante payment, as a function of type.\footnote{We did this by analytically solving the aggregator’s problem pointwise over \( \Omega_{DA} \), and then taking expectations of the quantities with respect to our market statistics.}

We also display as a dotted red line what the payment would be, if the aggregator maintained the same curtailment allocation, but were able to perfectly price discriminate (paying each increment its reservation price for curtailment, rather than the reservation price of the highest curtailed increment). The pink region between those two red lines is the “information rent”: the surplus payment that the increments are able to extract by virtue of their private information.

In Figure 4, we eliminate the type parameter from the contract menu, by plotting the payment to each increment type against the probability that that type is curtailed. This gives a more realistic depiction of a menu that the aggregator might offer in practice, assuming that actual participants have difficulty estimating the dollar value of each increment of their consumption. Here we see that the payment rises sub-linearly in probability. If we imagine only a single tuple of prices \( (p, a, b) \), then the probability of curtailment is piecewise linear: there is no curtailment for valuations above \( b/2 \), valuations below \( a/2 \) are always curtailed, and between those two levels, the probability varies linearly in type, as determined by the uniform distributions over wind and type (see Equation (15)). However, while the probability of curtailment falls piecewise linearly in type, the corresponding valuation for service rises linearly. If each increment were paid its reservation utility (i.e. no private information) the payment would be a concave quadratic function of type, and the lowest type would accept curtailment for no payment. With private information, the payment
curve has a similar quadratic form, but it is decreasing in type, and the payments are inflated. (Incentive compatibility, combined with merit order curtailment, forces the payment to be nonincreasing in type. Merit order curtailment implies that service plans designed for higher types have a lower probability of curtailment. There is no way to compel a low-type increment to accept a package intended for it, with high probability of curtailment and low payment, if an alternative package is available for higher types, which has lower probability of curtailment and a higher payment.) When the market prices are stochastic, the payment as a function of type is an average, or ex ante expectation, over such decreasing piecewise quadratic curves.

4 Discussion and conclusion

4.1 Under-purchase of demand response

We mentioned in Section 2 that the aggregator purchases less than the efficient level of demand response. This is because the aggregator is a monopsony purchaser of DR contracts. Under current institutional norms, customers effectively have option rights to consumption in the quantity of their physical fuse. But going forward, particularly with the rise of distributed generation, we anticipate that demand charges will become more prevalent. A demand charge requires the consumer to “buy the baseline,” or the capacity rights, which they would then re-sell to the aggregator. If the increments purchase less firm capacity up-front, then more flexibility will be available for real time adjustment and recourse in general. In order to investigate this, a model must incorporate an antecedent stage in which the DR participants purchase the baseline, as in, for example, Crampes and Léautier (Crampes and Léautier, 2015).
The incorporation of a demand charge would typically remove lower-type increments from the distribution (or demand curve) faced by the aggregator. The effect of this can be ambiguous for the aggregator’s profit. On the one hand, increments with low but positive net ex post valuation can provide cheap DR for the aggregator to purchase and deploy in the wholesale market; removing these from the market would reduce aggregator profit. On the other hand, increments with ex post valuation below the retail rate must be paid despite that they provide no DR, and act as a fixed cost for positive quantities of DR in each scenario. In our Example 2, if we add increments with valuation below the retail rate, the aggregator would always have to pay them to dispatch a positive quantity of DR, but it would not get demand response from them. The result is that in each RT scenario, the aggregator would purchase either the same amount of DR as if there were no such increments (because of marginal cost calculations), or it the aggregator would switch to purchasing zero DR in that scenario. We may explore the effects of demand charges more systematically in later work. Generally, we anticipate that a properly set demand charge can increase social welfare, but that a welfare-improving demand charge has an ambiguous effect on the aggregator’s profits, depending on specific conditions.

The other condition that results in under-purchasing of DR is the fact that the aggregator is a monopsonist, and consumers have private information. With either perfect competition or perfect information in the demand response contracting market, standard analysis shows that the quantity of demand response dispatched will be the ex post efficient quantity, assuming that imbalance prices are set competitively. A useful extension of our approach may be to model imperfect competition in the demand response market using residual demand curves.
4.2 Future work

In future work, we will demonstrate how to approximately solve the aggregator’s problem by simulation, from arbitrary distributions that satisfy our technical conditions.

We have seen the standard result, that the aggregator can do no better than quoting a single “clearing” price per scenario. This motivates an intuitive extension of our approach, abandoning the FOSD ordering assumption, and directly optimizing with respect to state-dependent demand response price bids.

5 Acknowledgment

This work was supported by the Department of Energy, through a grant administered by the Center for Electric Reliability Technology Solutions (CERTS).

A Appendix: Proofs

Proof of Lemma 2.8. Assumption 1 ensures the conditions sufficient for this claim. Note that $L(\cdot)$ is absolutely continuous, and differentiable except at two points. Further, recall that the support of $F(\cdot|\tau, \omega)$ is constant. Therefore we can apply integration by parts, and the antiderivative term drops out:

$$U(\hat{\tau}|\tau) = \int_{\Omega_{RT}} \int_{\Theta} -L'(\theta) F(\theta|\tau, \omega) d\theta k(\tau, \omega) dP_{RT}(\omega).$$

The partial derivative in question is therefore

$$\frac{\partial}{\partial \tau} U(\hat{\tau}|\tau) = \int \int -L'(\theta) \frac{\partial F(\theta|\tau, \omega)}{\partial \tau} k(\tau, \omega) dP_{RT}(\omega) \leq 0.$$

Each term in the integrand is uniformly bounded, which guarantees that this partial derivative is uniformly bounded, so that the function $U(\hat{\tau}|\cdot)$ is Lipschitz continuous, and thus absolutely continuous.

Proof of Proposition 2.10. By rearrangement, the payment is the equilibrium utility minus the net option value:

$$t(\tau) = u(\tau) - U(\tau).$$

First we simplify the expression for $u(\tau)$ from line (8). Under merit order curtailment, $k(\tau, \omega) = 1_{[\tau \leq \hat{\tau}(\omega)]}$. Exchanging the order of integration so that we integrate first with respect to $\tau$, we get that:

$$u(\tau) = u(\tau) - \int_{\Omega_{RT}} \int_{\Theta} L'(\theta)(f(\theta|\tau(\omega), \omega) - f(\theta|\tau, \omega)) k(\tau, \omega) dP_{RT}(\omega).$$

\[32\]

See Börgers (2010), end of chapter 2.
(a) Case 1: $\nu(b - a)/2 \leq 1$.

<table>
<thead>
<tr>
<th>Region $[q, \bar{q}_i]$</th>
<th>$[q, \bar{q}_i] =$</th>
<th>for $q \in [q, \bar{q}<em>i]$, $J</em>{DA}(q) = $</th>
<th>if $q_i &lt; q^* &lt; \bar{q}_i$, then $q^* = $</th>
</tr>
</thead>
<tbody>
<tr>
<td>1a</td>
<td>$0 &lt; q &lt; \nu a/2$</td>
<td>$(p - a)q + (a/2)(\nu a/2 + 1)$</td>
<td>$p &lt; a$</td>
</tr>
<tr>
<td>1b</td>
<td>$\nu a/2 &lt; q &lt; \nu b/2$</td>
<td>$(p - a - \frac{\nu^2 a}{4})q + \frac{\nu^2 a}{2}q^2 - \frac{1}{3\nu}q^3 + \frac{a}{2} + \frac{\nu a}{2} + \frac{\nu^2 a}{4}$</td>
<td>$a \leq p \leq a + \nu(\frac{a-b}{2})^2$</td>
</tr>
<tr>
<td>1c</td>
<td>$\nu b/2 &lt; q &lt; 1 + \nu a/2$</td>
<td>$(p - a - \frac{\nu^2 a}{4} + \frac{\nu a}{4})q + (\frac{a-b}{2})q^2 + \frac{\nu a}{24} + \frac{a}{2} - \frac{\nu a}{24} + \frac{\nu^2 a}{4}$</td>
<td>$a + \nu(\frac{a-b}{2})^2 \leq p \leq b - \nu(\frac{a-b}{2})^2$</td>
</tr>
<tr>
<td>1d</td>
<td>$1 + \nu a/2 &lt; q &lt; 1 + \nu b/2$</td>
<td>$(p + 2\nu a + \nu a)q - \left(\frac{b}{2} + \frac{1}{\nu}\right)q^2 + \frac{1}{3\nu}q^3 - \frac{1}{24}b^3 + \frac{1}{5\nu}$</td>
<td>$b - \nu(\frac{a-b}{2})^2 &lt; p &lt; b$</td>
</tr>
<tr>
<td>1e</td>
<td>$1 + \nu b/2 &lt; q$</td>
<td>$(p - b)q + \frac{b}{2} + \frac{\nu b}{4}$</td>
<td>$b &lt; p$</td>
</tr>
</tbody>
</table>

(b) Case 2: $\nu(b - a)/2 > 1$, $\iff 1 + \nu a/2 \leq \nu b/2$. Differences from Case 1

<table>
<thead>
<tr>
<th>Region $[q, \bar{q}_i]$</th>
<th>$[q, \bar{q}_i] =$</th>
<th>for $q \in [q, \bar{q}<em>i]$, $J</em>{DA}(q) = $</th>
<th>if $q_i &lt; q^* &lt; \bar{q}_i$, then $q^* = $</th>
</tr>
</thead>
<tbody>
<tr>
<td>2a</td>
<td>$0 &lt; q &lt; \nu a/2$</td>
<td>$(p - a)q + (a/2)(\nu a/2 + 1)$</td>
<td>$p &lt; a$</td>
</tr>
<tr>
<td>2b</td>
<td>$\nu a/2 &lt; q &lt; 1 + \nu a/2$</td>
<td>$(p - a - \frac{\nu^2 a}{4})q + \frac{\nu^2 a}{2}q^2 - \frac{1}{3\nu}q^3 + \frac{a}{2} + \frac{\nu a}{2} + \frac{\nu^2 a}{4}$</td>
<td>$a \leq p \leq a + 1/\nu$</td>
</tr>
<tr>
<td>2c</td>
<td>$1 + \nu a/2 &lt; q &lt; \nu b/2$</td>
<td>$(p + \frac{1}{\nu} + \frac{\nu a}{2})q - \frac{q^2}{\nu} - \frac{1}{3\nu}$</td>
<td>$a + 1/\nu \leq p \leq b - 1/\nu$</td>
</tr>
<tr>
<td>2d</td>
<td>$\nu b/2 &lt; q &lt; 1 + \nu b/2$</td>
<td>$(p + 2\nu a + \nu a)q - \left(\frac{b}{2} + \frac{1}{\nu}\right)q^2 + \frac{1}{3\nu}q^3 - \frac{1}{24}b^3 + \frac{1}{5\nu}$</td>
<td>$b - 1/\nu &lt; p &lt; b$</td>
</tr>
<tr>
<td>2e</td>
<td>$1 + \nu b/2 &lt; q$</td>
<td>$(p - b)q + \frac{b}{2} + \frac{\nu b}{4}$</td>
<td>$b &lt; p$</td>
</tr>
</tbody>
</table>

Table 1: $J_{DA}(q)$ and the optimal offer quantity $q^*$ in Example 2.
Proof of Proposition 2.12. Consider the difference in utility between the case where an increment of ex ante type $\tau$ reports truly, versus misreporting as $\tilde{\tau}$.

\[ u(\tau) - \left( U(\tilde{\tau}) + t(\tilde{\tau}) \right) \]

\[ = \int \int_{\Theta} L(\theta) f(\theta|\tau, \omega) d\theta \left( k(\tau, \omega) - k(\tilde{\tau}, \omega) \right) dP(\omega) + t(\tau) - t(\tilde{\tau}) \]

\[ = \int \int_{\Theta} L(\theta) \left( f(\theta|\tau, \omega) - f(\theta|\tilde{\tau}(\omega), \omega) \right) d\theta \left( k(\tau, \omega) - k(\tilde{\tau}, \omega) \right) dP(\omega) \]

\[ = \int \int_{\Theta} \left( -L(\theta) \frac{\partial}{\partial \tau} \left( F(\theta|\tau, \omega) - F(\theta|\tilde{\tau}(\omega), \omega) \right) \right) d\theta \left( k(\tau, \omega) - k(\tilde{\tau}, \omega) \right) dP(\omega). \]

Suppose $\tau > \tilde{\tau}$.Merit order curtailment implies that $k(\tau, \omega) - k(\tilde{\tau}, \omega) = -\mathbb{1}_{\{\Omega \leq \tilde{\tau}(\omega) < \tau\}} \leq 0$. On this set, FOSD implies that $F(\theta|\tau, \omega) \leq F(\theta|\tilde{\tau}(\omega), \omega)$.

This implies that the integral’s value is nonnegative. Similarly, if $\tau < \tilde{\tau}$, then $k(\tau, \omega) - k(\tilde{\tau}, \omega) = \mathbb{1}_{\{\tau \leq \tilde{\tau}(\omega) < \tilde{\tau}\}} \geq 0$. On this set, $F(\theta|\tau, \omega) \geq F(\theta|\tilde{\tau}(\omega), \omega)$. Again, the integral’s value is nonnegative.

Proof of Proposition 2.14. The substance of the Proposition is that (ignoring the penalty receipts term)

\[ \int_{\Omega} t(\tau) g(\tau) d\tau = -\int_{\Omega_R^T} \int_{\Omega} G(\tau) \frac{\partial}{\partial \tau} z(\tau, \omega_{RT}) + z(\tau, \omega_{RT}) g(\tau) d\tau. \]

(33)

Integrating Equation (9) under merit order curtailment, we get that

\[ \int_{\Omega} t(\tau) g(\tau) d\tau = \int_{\Omega_R^T} \int_{\Omega} z(\hat{\tau}(\omega_{RT})) g(\tau) dP(\omega_{RT}) \]

(34)

\[ = \int_{\Omega_R^T} z(\hat{\tau}(\omega_{RT})) G(\hat{\tau}(\omega_{RT})) dP(\omega_{RT}). \]

(35)

To obtain the desired formula, we differentiate the integrand in (35) using the product rule, factor out a “$g(\tau)$,” and re-integrate.

Proof of Proposition 3.8 (FONC for $q^\ast$). In order to differentiate this objective with respect to $q$, we decompose RT outcomes into three sets: overproduction $= \{DR + s > q\}$, shortfall $= \{DR + s < q\}$, and no imbalance $= \{DR + s = q\}$. From expression (15), we see that $\frac{\partial}{\partial q} \hat{q}^\ast(\omega_{RT}|q) = 1_{{\Omega(R)}}(\hat{\tau}(\omega_{RT}) = q - s)$. (Here

\[ \frac{1}{\delta} \frac{\partial}{\partial \omega_{DA}^T} \Omega(\hat{\tau} = \hat{\tau}^*|\omega_{DA}) = \frac{1}{\delta (\hat{\tau}) \hat{\tau}^*} \Omega(\hat{\tau} = \hat{\tau}^*|\omega_{DA}). \]

(36)

Despite the fact that the DR policy is held at its optimal values, we cannot apply an envelope theorem, because there is a positive probability that the value function is not differentiable with respect to $\tau$ at $\tau = \hat{\tau}^\ast$ (in all $\omega_{DA}^T$ events where there is zero imbalance at the optimum, the right derivative with respect to $\tau$ has an $a$ coefficient, and the left derivative has a $b$ coefficient).

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we ignore certain “edge” events, which we assume have probability zero:

abusing notation, $s = \inf\{s : DR(\hat{\tau}^*(s)) + s = q\}$, and $s = \sup\{s : DR(\hat{\tau}^*(s)) + s = q\}$. Also, note that the contributions to the expected derivative from the overproduction and shortfall terms are both zero when $DR + s = q$, because by assumption, the shortfall quantity is constantly zero on this set. When overproduction and underproduction are strict, then $\frac{\partial}{\partial q} \hat{\tau}^*(\omega_{RT}|q) = 0$.

This gives us that

$$
\frac{\partial}{\partial q} J(q|\omega_{DA}) = p - \mathbb{E}[a|DR(\hat{\tau}^*) + s > q] \Pr(DR(\hat{\tau}^*) + s > q) - \mathbb{E}[b|DR(\hat{\tau}^*) + s < q] \Pr(DR(\hat{\tau}^*) + s < q) - \mathbb{E}[c(\hat{\tau}^*)|DR(\hat{\tau}^*) + s = q] \Pr(DR^* + s = q).
$$

The first order condition follows.

References


